Bolt on some Crypto

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Securing The Network - TLS & SSH

IETF Standards:

SSH - RFC 4250-4255

- > Remote shell
- > File transfer
- > TCP port forwarding, socks proxy
- Pipe commands over ssh (stdin/stdout)
- Originally a replacement for BSD r-commands

TLS - RFC 5246.

- ➤ https://
- Optional for SMTP, IMAP, POP3, XMPP, LDAP

TLS & SSH - Cryptographic Services

- Authentication
 - more to come on this...
- > Integrity
 - Any tampering with the connection will be detected
 - Limitation: attacker can drop the session
 - Limitation: DoS
- > Privacy
 - Cannot see contents of session
 - Limitation: traffic analysis (aka metadata)

Public Key Cryptography Primer

Keypair:

Private Key - This key must be kept safe! *Don't email me your private key!*

Public Key - This key can be shared with anyone you need to communicate with

Signing:

The Private Key is used to sign a hash of a message, which can be verified by anyone with the public key

Encryption:

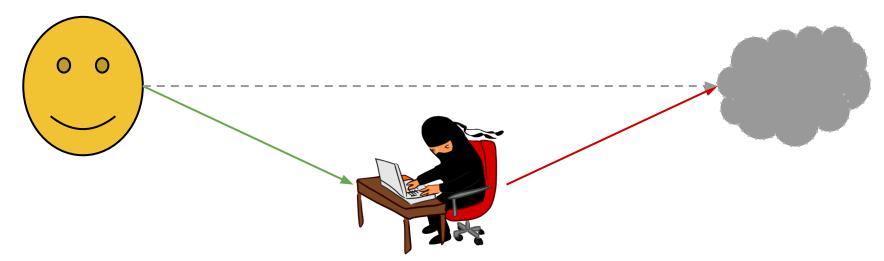
The Public Key is used to encrypt a message, which only the holder of the Private Key can decrypt

MiTM Attack

"Man-in-the-middle attack"

(The actual attack isn't gender specific)

- 1. Intercept client connection and answer like a server.
- 2. Connect to the real server (optional)
- 3. Log or modify data as it passes through



MiTM Attack - Linux Quickstart

- iptables -t nat -A PREROUTING -p tcp --dport 5222 -j \
 REDIRECT --to-port 5002
- Run your client program, listening on 5002
- > Route the traffic *through* your linux box using arpspoof, routing protocol
- If using dns spoofing, IP tables not required

To get the original dest IP:

In C:

```
getsockopt(s, SOL_IP, SO_ORIGINALDEST, &addr, &addrlen);
In Python:
packedDest = s.getsockopt(socket.SOL_IP, 80, 16)
(destPort, ) = struct.unpack(">H", packedDest[2:4])
destHost = socket.inet ntoa(packedDest[4:8])
```

SSH Host Keys

OpenSSH caches host keys:

StrictHostKeyChecking - on seeing a new host:

- > yes don't connect
- no cache host key and connect
- ask (default) display the fingerprint and ask user

ssh-keyscan can collect host keys from remote systems - allows you to pre-populate known_hosts

SSH Host Keys - API gotchas

JSch:

StrictHostKeyChecking=no won't cache the host key!

Paramiko:

```
client.load_system_host_keys()
client.set_missing_host_key_policy(paramiko.RejectPolicy)
```

paramiko.WarningPolicy won't cache the host key!

Just pre-populate /etc/ssh/ssh_known_hosts if using APIs - no need for write access to known_hosts

SSH MiTM - you can do 'em

Often StrictHostKeyChecking=no is set on servers with unattended ssh sessions

- > rsync jobs
- > remote commands

OpenSSH still connects if the host key changed and you're using public key authentication!

A MiTM server could just accept pubkey auth for any key (without knowing the key).

SSH Client Authentication

You can create a client keypair with ssh-keygen, then add it to ~/.ssh/authorized_keys on remote hosts. This can be put in kickstart/preseed files.

Even if the remote server is compromised your private key should be safe, so you don't need a fresh one for each server you connect to.

You can do "two-factor" in OpenSSH with the AuthenticationMethods sshd_config option.

TLS - X.509 Certificates

Certificate chain from my website:

```
0 s:/CN=www.miknet.net
i:/CN=StartCom Class 1 Primary Intermediate Server CA
1 s:/CN=StartCom Class 1 Primary Intermediate Server CA
i:/CN=StartCom Certification Authority
```

StartCom Certification Authority is trusted by my system

Subject: the entity identified by the certificate

Issuer: the authority that signed the certificate

Domain Validated: demonstrated control of the domain to CA

Extended Validation: demonstrated that you are the organisation and domain holder in the certificate

The unverified certificate

A Root CA is just a self-signed certificate Intermediate CAs and the certificate are signed by their parent CA

You can create an entire *unverified* chain using the

openssl command line. Only the

public key matters.

Even the most diligent support staff would tell users to click through.



Dialogs that shouldn't exist





WARNING: TLS APIs suck

There are 3 types of TLS APIs:

- Go verify the certificate yourself
 - Generally OpenSSL or wrappers
- What's a certificate?
 - High level abstractions over OpenSSL written by programmers who don't know/understand
 - all of the Python 2.x standard library
- We do what a web browser would
 - These are rare python-requests.org, libcurl

TLS - Verifying the hostname

Most TLS libraries do not check that the certificate matches the hostname - even if you turn on verification.

Should you trust a certificate for www.miknet.net when accessing your online banking?

The hostname must match either the **CN** field or one of the **SubjectAltName** extensions.

WARNING: NULL bytes are valid

Match the name the user requested, not DNS SRV/MX

STARTTLS

```
<?xml version='1.0' ?>
<stream:stream to='jabber.org' xmlns='jabber:client' xmlns:stream='http:</pre>
//etherx.jabber.org/streams' version='1.0'>
<?xml version='1.0'?>
<stream:stream xmlns='jabber:client' xmlns:stream='http://etherx.jabber.</pre>
org/streams' from='jabber.org' id='5ce74cfce8e91fc4' version='1.0'>
  <stream:features>
    <starttls xmlns='urn:ietf:params:xml:ns:xmpp-tls'/>
    <mechanisms xmlns='urn:ietf:params:xml:ns:xmpp-sasl'>
      <mechanism>DIGEST-MD5</mechanism>
      <mechanism>PLAIN</mechanism>
    </mechanisms>
  </stream:features>
  <starttls xmlns='urn:ietf:params:xml:ns:xmpp-tls'/>
  ceed xmlns='urn:ietf:params:xml:ns:xmpp-tls'/>
```

STARTTLS removed

```
<?xml version='1.0' ?>
<stream:stream to='jabber.org' xmlns='jabber:client' xmlns:stream='http:</pre>
//etherx.jabber.org/streams' version='1.0'>
<?xml version='1.0'?>
<stream:stream xmlns='jabber:client' xmlns:stream='http://etherx.jabber.</pre>
org/streams' from='jabber.org' id='5ce74cfce8e91fc4' version='1.0'>
  <stream:features>
    <starttls xmlns='urn:ietf:params:xml:ns:xmpp-tls'/>
    <mechanisms xmlns='urn:ietf:params:xml:ns:xmpp-sasl'>
      <mechanism>DIGEST-MD5</mechanism>
      <mechanism>PLAIN</mechanism>
    </mechanisms>
  </stream:features>
  <starttls xmlns='urn:ietf:params:xml:ns:xmpp-tls'/>
  ceed xmlns='urn:ietf:params:xml:ns:xmpp-tls'/>
```

STARTTLS removal

Don't negotiate whether to encrypt *over the network*!

- > XMPP, IMAP, POP3, SMTP clients
- SMTP server-to-server always works
- Dell & Cisco BMCs that use the Avocent KVM stack (PoC||GTFO 0x5)
- > HTTP (sslstrip by Moxie)

Forward Secrecy

Ephemeral Key Exchange is another form of public key cryptography

- > Protocols: Diffie-Hellman or Elliptic Curve Diffie-Hellman
- Known as: Forward Secrecy or PFS
- > TLS Ciphersuits that start with DHE- or ECDHE-
- The SSLv3 ciphersuites use RSA encryption if the RSA key is stolen/cracked, past traffic can be decrypted! (Wireshark supports this)

Forward Secrecy - TLS Ciphersuites

Apache:

```
SSLCipherSuite ...

SSLHonorCipherOrder on

SSLProtocol all -SSLv2 -SSLv3
```

Nginx:

```
ssl_protocols TLSv1 TLSv1.1 TLSv1.2;
ssl_prefer_server_ciphers on;
ssl_ciphers "...";
```

5 Ciphers that have you covered (thanks to Kenn White):

```
ECDHE-RSA-AES256-GCM-SHA384 # Android 4.4+

ECDHE-RSA-AES128-SHA256 # IE 11

ECDHE-RSA-AES128-SHA # Android 4.x, Chrome/Firefox, IE8-10

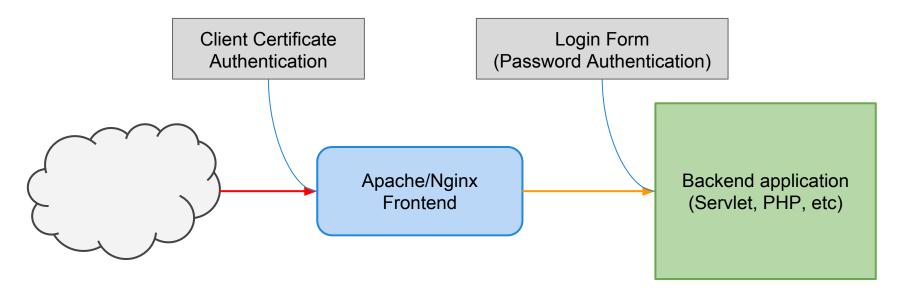
DHE-RSA-AES128-SHA # Android 2

RC4-SHA # Old junk (Windows XP, Nokia 6xxx) RC4 MUST DIE!
```

TLS - Authenticating Clients

Originally SSL was for e-commerce. This only required "money green" authenticity for clients.

TLS has support for client certificates



Entropy

Most cryptography needs randomness, for both short-term and long term keys.

The properties that are needed:

- Unable to predict future values
- Unable to recover past values

PRNGs work but need to be seeded from truly unguessable events.

Not Entropy

- mt_rand()
 - Can recover all state from output.
 - Often a small input
- rand() / random()
 - Small input
 - Can recover some/all state from output
- rand_r() / qrand() / java.util.Random
 - Small input
 - Small state
 - Can recover some state from output

https://www.miknet.net/rux2013/

Entropy - Don't fork it up

- Unix-like systems: read from /dev/urandom
 - Userland PRNGs probably not fork() safe
- Windows: CryptGenRandom for strong entropy

- Linux early boot (only): /dev/random
 - Encrypted swap
 - SSH host key generation

Hash Functions - attack types

A fixed-length digest of variable length input

- > (First)-Preimage resistance
 - Hard to find the original input from the hash
 - Guessing inputs still works!
- Second-Preimage resistance
 - Hard to find a second input that produces a given hash
 - An ideal hash function would provide 2^{hash length} resistance to this
- Collision Resistance
 - Hard to find two inputs that produce the same hash
 - Birthday attack requires 256-bit hash for 128-bit security
 - When a hash function is broken this is usually first to go

Hash Functions - attack examples

- If the attack controls multiple inputs, you need to worry about collisions
 - rsync/librsync (see my github)
 - X.509 (TLS) certificates
- If the attacker controls one input, you need to worry about second-preimage

Checksums and Signatures

For general use, you should use SHA-2 256/384/512 Creating a certificate:

openssl req -new -sha256 ...

Checksum of a file:

sha256sum *.iso

blake2 - is a very fast and secure hash function - if performance is critical.

Don't use MD4/MD5 at all

SHA-1 should be phased out

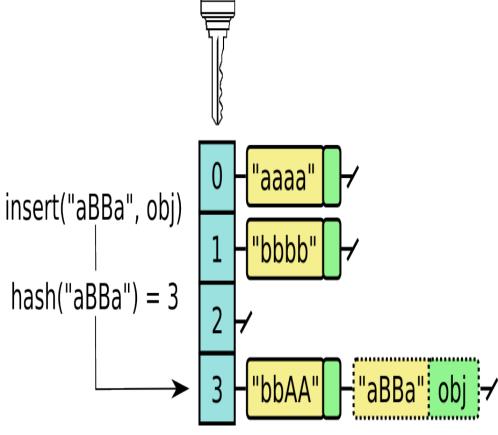
Hash Tables

Wait, do these need to be secure?!?

Worker pools or select() loops - colliding hash table entries can block the CPU!

SipHash was designed to fix this

This is now the default in Python3, Ruby, Perl



http://commons.wikimedia.org/wiki/File:Hashtable_linkedlist_collision.png

Password Hashes

Normal hash functions allow you to make extremely fast guesses - do not use these!

Salting

A salt is a unique string that is hashed with the password and stored next to the hash.

- Mutliple users with the same password won't have the same hash
- An attacker can't pre-calculate passwords

Stretching

- An operation that makes the hashing deliberately slow
- Must be sure that attackers can't take a shortcut

Current recommendation: bcrypt

Future recommendation: Winner of PHC - https://password-hashing.net/

Advanced Password Hashing - HSM

Passwords are still really weak! Solar Designer: encrypt hashes with a HSM

- Near-perfect security if HSM is safe
- Better than nothing if HSM is stolen/broken
- Only requires encrypt function of HSM

Store:

```
salt, AES(bcrypt(password, salt))
Compare:
AES(bcrypt(password, salt)) == stored
```

http://www.openwall.com/presentations/YaC2012-Password-Hashing-At-Scale/http://www.openwall.com/presentations/Passwords12-The-Future-Of-Hashing/

Password Hashing - Remediation

If you store your passwords in cleartext - go hash them all!

If you use an unsalted hash - use a password hash on the original hash

Avoid HTTP Digest Authentication, NTLM, CHAP.

MACs

If you need an untrusted entity to hold some state for you, you can use a MAC

- > Ensure your data cannot be used out of context
 - HKDF, or just separate keys
- > The key needs to be *secret* and preferably *random*
- ➤ Timing attacks! Brad Hill's trick:

```
HMAC(random, mac) == HMAC(random, HMAC(secret, data))
```

If the data needs to be encrypted MAC the ciphertext (EtM)

Hash functions aren't MACs!

Most hash functions do not function as a MAC.

Bad:

```
tag = Hash(secret, message)
and send tag, message that person can then
perform a length extension attack.
```

Use HMAC when there's a secret key

Homework

- Write a MiTM attack for TLS and/or SSH
- Try it against every connection that leaves your machine
- > File bug reports
- Code a length-extension attack
- cryptopals.com (if you like this stuff)

Thanks!

Comments/Questions?

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